

Preparing for Olympiads

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Outline

- 1 Introduction
- 2 Preparation
 - Training
 - Practise
 - Skills
 - Strategy
- 3 Summary

Who Am I?

- South African competition programmer
- IOI 2006–2011
- ACM ICPC 2002–2003
- USACO, Topcoder, Code Jam, Challenge24, IPSC, ICFP,
...

Who Am I?

- Post-doctoral researcher at UCT
- Research in GPU computing



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Why Am I Here?

Three talks today

- How I prepare for olympiads
- Turing Machines
- Geometry with complex numbers

Scope

I'll talk mostly about ICPC/Code Jam/TC Algorithm etc style:

- Computational problems
- Well-defined bounds on problem size
- Answers are either right or wrong
- Big-O complexity is important
- Micro-optimisation **not** important
- Objective testing using test data

The Secret To Success

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There isn't one.

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There isn't one.

- Natural ability
- Experience
- Training
- Practise
- Skills
- Strategy

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Algorithms

- Bellman-Ford
- Biconnected components
- Binary search
- Breadth-first search
- Bresenham's algorithm
- Delaunay triangulation
- Depth-first search
- Boyer-Moore
- Dijkstra's algorithm
- Dynamic programming
- Euclidean algorithm
- Eulerian paths
- Finite state machines
- Fast exponentiation
- Floyd-Warshall
- Ford-Fulkerson
- Graham scan
- Knuth-Morris-Pratt
- Linear programming
- Matching
- Min-cost network flow
- Minimax
- Memoisation
- Prim's algorithm
- Rabin-Karp
- Sorting
- Stable marriage problem
- Strongly connected components
- Recursion
- Rotating calipers
- Sorting algorithms
- Kruskal's algorithm

Data Structures

- Arrays
- Balanced binary trees
- Binary heaps
- Binary indexed trees
- Binary trees
- Bit vectors
- Deques
- Hash tables
- Interval trees
- Graphs
- Linked lists
- Queues
- Range trees
- Stacks
- Suffix trees
- Tries
- Union-find trees

Mathematics

- Proof techniques
- Complexity analysis
- NP-completeness
- Problem transformations
- Binary representations
- Combinatorics
- Modular arithmetic
- Invariants
- Probabilities
- Prime factorisations
- Linear algebra

Tips

To master an algorithm/structure

- Learn it

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- Implement it

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- Teach it to someone else

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Tips

To master an algorithm/structure

- Learn it
- Implement it
- Teach it to someone else
- Use it until you've **memorised** it
- Be able to implement it **fast**

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Do Contests

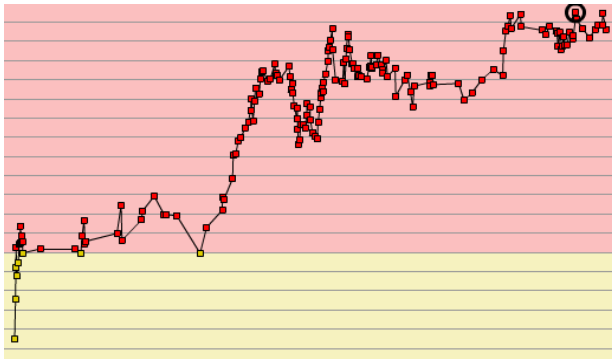
Top 10 Topcoder Algorithm competitors:

Rank	Handle	Matches
1	Petr	288
2	tourist	111
3	ACRush	178
4	bmerry	170
5	nika	139
6	tomek	171
7	wata	174
8	UdH-WiNGeR	96
9	marek.cygan	219
9	dzhulgakov	96

As of 15/04/2012

Do Contests

My Topcoder rating history



Learn From Mistakes

- Debug submissions that failed

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- Debug submissions that failed
- Learn how to solve the problems you couldn't

Learn From Mistakes

- Debug submissions that failed
- Learn how to solve the problems you couldn't
- See how others solved it

Adapt To Avoid Mistakes

Binary Search

- Find largest i in $[0, N - 1]$ s.t. $f(i)$ is true
- $f(0)$ is true, $f(i + 1) \Rightarrow f(i)$

Adapt To Avoid Mistakes

Binary Search

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```
int L = 0; // search [L, R]
int R = N - 1;
while (L != R) {
    int M = (L + R) / 2;
    if (f(M))
        L = M;
    else
        R = M - 1;
}
return L;
```

Adapt To Avoid Mistakes

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```

```
int L = 0; // f(L) is true
int R = N; // f(R) is false
while (R - L > 1) {
    int M = (L + R) / 2;
    if (f(M))
        L = M;
    else
        R = M;
}
return L;
```

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Tools

Software tools can find errors

- Compiler flags
- Debugger
- Profiler
- Memory checker
- Library assertions

Libraries

Don't reinvent the wheel

- Know the language containers in detail
- Know the complexity of operations

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Keep it simple

Everything should be made as simple as possible, but no simpler. — Albert Einstein (paraphrased)

Debugging is twice as hard as writing the code in the first place. Therefore, if you write the code as cleverly as possible, you are, by definition, not smart enough to debug it. — Brian W. Kernighan

Simplification Examples

- Use the simplest algorithm that is fast enough.

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- If in doubt, use 64-bit integers.
- Allocate slightly more memory than you need.
- Don't try to put everything on one line.
- Don't repeat yourself. Write a function.

Write Less Code

Part of a solution to a TC problem

```
public class EvenRoute {
    public String isItPossible(int[] x, int[] y, int wantedParity) {
        boolean hasFP = false, hasFU = false;
        for (int i=0; i<x.length; ++i)
            if ((1 & (x[i]+y[i])) == 0)
                hasFP = true; else hasFU = true;
        boolean hasN = false;
        for (int i=0; i<x.length; ++i) for (int j=0; j<x.length; ++j)
            if (((x[i]-x[j] + y[i]-y[j]) & 1) == 1) hasN = true;
            if (wantedParity == 0) {
                if (!hasFP && !hasN) return "CANNOT";
                else return "CAN";
            } else {
                if (hasFU || hasN) return "CAN";
                else return "CANNOT";
            }
        }
    }
}
```

Write Less Code

Another solution:

```
public class EvenRoute {  
    public String isItPossible(int[] x, int[] y, int wantedParity) {  
        int n = x.length;  
        for (int i = 0; i < n; ++i) {  
            if (Math.abs(x[i] + y[i]) % 2 == wantedParity) {  
                return "CAN";  
            }  
        }  
        return "CANNOT";  
    }  
}
```

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- If it seems too easy, it might be

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- If it seems too easy, it might be
- If you can't prove your algorithm is correct, try to break it
- Use assertions in your code
- **Assume** your code has a bug

Summary

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- Learn how to solve problems

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- Keep things simple
- Avoid subtle errors

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- Learn how to solve problems
- Keep things simple
- Avoid subtle errors
- **Have fun**